

**MATRIX RATIO TEST FOR FINDING BASIC
FEASIBLE SOLUTION WITHOUT ARTIFICIAL
VARIABLE(S) IN LP PROBLEMS**

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Abstract: Note, that each artificial variable creates one extra dimension for LP problem. In this paper, we show how can find Basic Feasible Solution *without adding artificial variable(s)*.

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1. Introduction

Consider the following LP problem:

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$$\begin{aligned} \text{Min} \quad & cx \\ \text{s.t.} \quad & Ax = b \\ & x \geq 0 \end{aligned}$$

that A is the $m \times n$ Matrix and $\text{Rank}(A) = m$.

Suppose that we can partition the constraint matrix A into $A = [B, N]$. Where B is a Basic matrix, not necessarily feasible.

We known three methods to initiate the simplex method algorithms, by the use of artificial variable(s):

- i) Big-M method [3, Section 2.7].
- ii) Two phase Method [3, Section 2.4].
- iii) Single variable technique [1, Section 4.4], [2, Section 4.5] and [2, Exercise 4.33].

We want to show in many problems, *with our technique*, we can find *basic feasible solution*, without using artificial variable(s).

2. Maximum Ratio Test (MRT)

The goal is to find *basic feasible solution*. Without adding artificial variable(s). Consider following tableau:

	x_k	RHS
	y_{1k}	\bar{b}_1
.	.	.
.	.	.
x_{B_p}	y_{pk}	\bar{b}_p
.	.	.
.	.	.
.	y_{rk}	\bar{b}_r
.	y_{rk}	\bar{b}_r
.	$y_{(r+1)k}$	\bar{b}_{r+1}
.	.	.
.	.	.
.	y_{mk}	\bar{b}_m

Suppose $\bar{b} \not\geq 0$:

If $\bar{b}_i < 0$ then $i \in J_1 = \{1, \dots, r\}$.

If $\bar{b}_i \geq 0$ then $i \in J_2 = \{r + 1, \dots, m\}$.

In tableau, if there was y_k subject to $y_{ik} < 0$ for $i \in J_1$ and $y_{ik} \leq 0$ for $i \in J_2$. We can use MRT to find a basic feasible solution:

$$\frac{\bar{b}_p}{y_{pk}} = \text{Maximum} \left\{ \frac{\bar{b}_i}{y_{ik}}, i \in J_1 \right\}.$$

3. Validation of MRT

In previous tableau, if x_k enters the basic and x_{B_p} , leaves the basic, then pivoting can be done on y_{pk} :

	x_k	RHS
.	0	.
.	.	.
x_k	1	$\frac{\bar{b}_p}{y_{pk}}$
.	.	.
.	.	.
x_{B_i}	.	$\bar{b}_i - \frac{\bar{b}_p}{y_{pk}}y_{ik}$
.	.	.
.	0	.

We show that after pivoting, right hand side becomes positive:

- 1) $p \in J_1 \implies \frac{\bar{b}_p}{y_{pk}} > 0$.
- 2) $p \in J_1 \ \& \ i \in J_1 \implies \bar{b}_i - \frac{\bar{b}_p}{y_{pk}}y_{ik} = y_{ik}\left(\frac{\bar{b}_i}{y_{ik}} - \frac{\bar{b}_p}{y_{pk}}\right) \geq 0$.
- 3) $p \in J_1 \ \& \ i \in J_2 \implies \bar{b}_i - \frac{\bar{b}_p}{y_{pk}}y_{ik} = y_{ik}\left(\frac{\bar{b}_i}{y_{ik}} - \frac{\bar{b}_p}{y_{pk}}\right) \geq 0$.

Corollary. Each basic solution that has got this y_k , its problem is feasible.

4. Numerical Example

Example. Solve the following problem:

$$\begin{aligned} &\text{Maximum} && 2x_2 - x_2 + x_3, \\ &s.t. && x_1 + x_2 - x_3 \geq 4, \\ &&& x_1 - 4x_2 + x_3 \leq 2, \\ &&& x_1, x_2, x_3 \geq 0. \end{aligned}$$

$$\begin{aligned}
 \text{Maximum} \quad & 2x_1 - x_2 + x_3, \\
 \text{s.t.} \quad & -x_1 - x_2 + x_3 \leq -4, \\
 & x_1 - 4x_2 + x_3 \leq 2, \\
 & x_1, x_2, x_3 \geq 0.
 \end{aligned}$$

	x_1	x_2	x_3	x_4	x_5	RHS
z	-2	1	-1	0	0	0
x_4	-1	-1	1	1	0	-4
x_5	1	-4	1	0	1	2

We enter x_2 in to the basic and with MRT x_4 leave the basic.

	x_1	x_2	x_3	x_4	x_5	RHS
z	-3	0	0	1	0	-4
x_2	1	1	-1	-1	0	4
x_5	5	0	-3	-4	1	18

RHS becomes positive, and we can solve the problem with simplex method.

5. Conclusion

In many basic solutions, we do not have such a y_k , that defined in Section 2, and our technique can be used just for basic solutions that have such a y_k .

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