

GOPPA CODES ON CURVES ON A HIRZEBRUCH SURFACE

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Abstract: Here we study Goppa codes on smooth curves contained in a Hirzebruch surface.

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1. Introduction

We work over a finite field K . We consider curves on a Hirzebruch surface and some of the Goppa codes constructed using them (in particular the minimum distance and a description of all the codewords with minimum weight). Let C be a smooth and geometrically connected curve defined over K . For any line bundle \mathcal{A} on C defined over K and any $B \subseteq C(K)$ let $\mathcal{C}(B, \mathcal{A})$ denote the code obtained evaluating $H^0(C, \mathcal{A})$ at the points of B ; if $\mathcal{A} \cong \mathcal{O}_C(D)$ with D an effective divisor of C defined over K and whose support contains no point of B , then $\mathcal{C}(B, \mathcal{A}) \setminus \{0\}$ is the set of all rational functions $f \in K(C)$ defined over K and with $(f) + D \geq 0$, ([2], Ch. 2, [3]) (it is the geometric Goppa code $C_{\mathcal{L}}(B, D)$ defined in [2], II.2.1). The dual code $\mathcal{C}(B, \mathcal{O}_C(D))^{\perp}$ may be described in the same way ([2], Theorem II.2.8). For any $\mathbf{w} \in \mathcal{C}(B, \mathcal{A})^{\perp} \setminus \{0\}$ let $\text{Supp}(\mathbf{w}) \subseteq B$ denote the support of \mathbf{w} . Hence the Hamming weight of \mathbf{w} is the integer $\#(\text{Supp}(\mathbf{w}))$.

For any integer $e > 0$ let $X_e := \mathbb{P}(\mathcal{O}_{\mathbb{P}^1} \oplus \mathcal{O}_{\mathbb{P}^1}(e))$ denote the Hirzebruch surface, i.e. the only ruled surface of \mathbb{P}^1 with a section with self-intersection $-e$ ([1], §V.2). X_e is a smooth and geometrically connected projective surface defined over K . Let $\pi : X_e \rightarrow \mathbb{P}^1$ be the ruling of X_e . We call $f = \pi^{-1}(P)$, $P \in \mathbb{P}^1(K)$, the class in $\text{Pic}(X_e)(K)$ of a fiber of π and as h both the class of a section of π with self-intersection $-e$ and the corresponding curve $h \subset X_e$ with $h \cong \mathbb{P}^1$. Since we assumed $e > 0$, the curve h is unique and it is defined over K . We have $f \cong \mathbb{P}^1$ and $h \cong \mathbb{P}^1$ and $\text{Pic}(X_e)(\overline{K}) \cong \text{Pic}(X_e)$. We have $h^2 = -2$, $h \cdot f = 0$ and $f^2 = 0$ (intersection numbers in $\text{Pic}(X_e)$).

Theorem 1. *Fix integers e, a, b, x, y such that $e > 0$, $b \geq ea > 0$, a smooth curve $C \in |\mathcal{O}_{X_e}(ah+bf)|$ defined over K and a zero-dimensional scheme $E \subset C$. Fix integers $x \geq ey > 0$ and a set $B \subset C(K) \setminus E_{reg}$ such that $\#(B) + \text{deg}(E) > ax + by - eay$. Assume either $y = a - 1$ or $y \geq a$ and $x - b \geq e(y - a) - 1$. Set $\mathcal{C}(B, \mathcal{O}_C(yh + xf)(-E))$, $m := \text{deg}(E \cap h)$ and $m' := \#(B \cap h)$. Let s be the maximal integer $\text{deg}(D \cap E)$ with $D \in |\mathcal{O}_{X_e}(f)|(\overline{K})$ and σ the set of all $D \in |\mathcal{O}_{X_e}(f)|(K)$ such that $\text{deg}(D \cap S) = s$. Let σ' be the set of all $D \in \sigma$ such that $D \cap E = \emptyset$. Then \mathcal{C}^\perp has minimum distance $\geq \min\{y + 2 - s, x - ey + 2 - m\}$. Let \mathcal{A} be the set of all $S \subseteq B$ such that $S \subset D$ for some $D \in \sigma$ and $\#(S) = y + 2 - s$. Let \mathcal{S} be the set of all $S \subseteq h \cap B$ such that $\#(S) = x - ey + 2 - m$ ($S = \emptyset$ if and only if $m' \leq x - ey + 1 - m$).*

(i) *If $y - s \leq x - ey - m$, then \mathcal{C}^\perp has minimum distance $y + 2 - s$ if and only if $\mathcal{A} \neq \emptyset$. If $\mathcal{A} \neq \emptyset$, then \mathcal{C}^\perp has $(\#(K) - 1)\#(\mathcal{A})$ minimum weight codewords (each codeword is uniquely determined, up to a non-zero constant, by its support and the supports are exactly the elements of \mathcal{A}).*

(ii) *If $y - s > x - ey - m$, then \mathcal{C}^\perp has minimum distance $x - ey + 2 - m$ if and only if $\mathcal{S} \neq \emptyset$. If $\mathcal{S} \neq \emptyset$, then \mathcal{C}^\perp has $(\#(K) - 1)\#(\mathcal{S})$ minimum weight codewords (each codeword is uniquely determined, up to a non-zero constant, by its support and the supports are exactly the elements of \mathcal{S}).*

(iii) *If $y - s = x - ey - m$, then \mathcal{C}^\perp has minimum distance $y + 2 - s$ if and only if $\mathcal{A} \cup \mathcal{S} \neq \emptyset$, \mathcal{C}^\perp has $(\#(K) - 1)\#(\mathcal{S} \cup \mathcal{A})$ minimum weight codewords and these codewords are described as in (i) and (ii).*

For the other parameters of the code \mathcal{C} , see Remark 3.

Remark 1. Take the set-up of Theorem 1, but drop the assumption that either $y = a - 1$ or $y \geq a$ and $x - b \geq e(y - a) - 1$. Then \mathcal{C}^\perp has minimum distance at least the one described in Theorem 1 and if equality holds, then the minimum weight codewords are described as in Theorem 1. Let $V \subset H^0(C, \mathcal{O}_C(yh + xf))$ be the image of the restriction map $H^0(C, \mathcal{O}_{X_e}(yh + xf)) \rightarrow H^0(C, \mathcal{O}_C(yh +$

xf). The statement of Theorem 1 applies to the code obtained evaluating V at the points of B .

2. The Proof

Let \overline{K} denote the algebraic closure of K . Let X be any variety or scheme defined over K and \mathcal{F} any coherent sheaf on X . Let $X_{\overline{K}}$ and $\mathcal{F}_{\overline{K}}$ denote the objects obtained extending the field of coefficients from K to \overline{K} . For each $i \in \mathbb{N}$ we have $\dim_K(H^i(X, \mathcal{F})) = \dim_{\overline{K}}H^i(X_{\overline{K}}, \mathcal{F}_{\overline{K}})$ ([1], Proposition III.9.3). Set $h^i(X, \mathcal{F}) := \dim_K H^i(X, \mathcal{F})$. This observation allows us to freely use un the next two remarks many results from [1], §V.2, only stated over \overline{K} .

Remark 2. Fix integers x, y . We have $h^0(X_e, \mathcal{O}_{X_e}(yh + xf)) = 0$ if and only if either $y < 0$ or $x < 0$. For all $y \geq -1$ we have $R^i\pi_*(\mathcal{O}_{X_e}(yh + xf)) = 0$ for all $i > 0$. For all $y \geq 0$ we have $\pi_*(\mathcal{O}_{X_e}(yh + xf)) \cong \bigoplus_{i=0}^y \mathcal{O}_{\mathbb{P}^1}(x - ie)$. Hence $h^i(X_e, \mathcal{O}_{X_e}(-h + xf)) = 0$, $i = 0, 1$, $h^1(X_e, \mathcal{O}_{X_e}(yh + xf)) = 0$ if $y \geq 0$ and $x \geq -1 + ey$ and $h^0(X_e, \mathcal{O}_{X_e}(yh + xf)) = \sum_{i=0}^y (x + 1 - ey) = (y + 1)(x + 1) - ey(y + 1)/2$ if $y \geq 0$ and $x \geq -1$. The line bundle $\mathcal{O}_{X_e}(yh + xf)$ is spanned if and only if $y \geq 0$ and $x \geq ey$ (if $y > 0$ and $0 \leq x < ey$, then h appears in the base locus of $|\mathcal{O}_{X_e}(yh + xf)|$ with multiplicity $\lceil (ye - x)/e \rceil$). For all integers $y > 0$ and $x \geq ey$ there is a smooth and geometrically connected curve $C \in |\mathcal{O}_{X_e}(yh + xf)|$ defined over K (use [1], Proposition 2.20). Set $g(C) := p_a(C)$. We have $\omega_{X_e} \cong \mathcal{O}_{X_e}(-2h - (e + 2)f)$ ([1], Lemma V.2.10). Hence the adjunction formula gives $\omega_C \cong \mathcal{O}_C((y - 2)h + (x - e - 2)f)$ ([1], Proposition II.8.20 or Proposition V.1.5). Hence $2g(C) - 2 = \mathcal{O}_{X_e}((y - 2)h + (x - e - 2)f) \cdot \mathcal{O}_{X_e}(yh + xf) = -ey(y - 2) + x(y - 2) + (x - e - 2)y$. Hence $g(C) = 1 + xy - x - y - ey(y - 1)/2$.

Remark 3. Fix integers x, y, a, b such that either $y = a - 1$ or $x - b \geq e(y - a) - 1$ and $b \geq ae > 0$. We have $h^1(X_e, \mathcal{O}_{X_e}((y - a)h + (x - b)f)) = 0$. Hence for each $C \in |\mathcal{O}_{X_e}(ah + bf)|$ the restriction map $H^0(C, \mathcal{O}_{X_e}(yh + xf)) \rightarrow H^0(C, \mathcal{O}_C(yh + xf))$ is surjective. Now assume $x \geq ey > 0$. We get $h^0(C, \mathcal{O}_C(yh + xf)) = h^0(X_e, \mathcal{O}_{X_e}(yh + xf)) - h^0(X_e, \mathcal{O}_{X_e}((y - a)h + (x - b)f)) = (y + 1)(x + 1) - ey(y + 1)/2 - (y + 1 - a)(x + 1 - b) - e(y - a)(y - a + 1)$ (Remark 2). Let $E \subset C$ be a zero-dimensional scheme. We have $\deg(\mathcal{O}_C(yh + xf)(-E)) = ay + bx - eay - \deg(E)$. If $h^1(X_e, \mathcal{I}_E(yh + xf)) = 0$, then $k := h^0(C, \mathcal{O}_C(yh + xf)(-E)) = h^0(C, \mathcal{O}_C(yh + xf)) - \deg(E)$. If $\sharp(B) > ay + bx - eay - \deg(E)$, then \mathcal{C} is an $[n, k]$ -code.

Remark 4. Let W be any projective scheme and L a line bundle on it. Fix any subscheme $E \subseteq Z$. Since Z is zero-dimensional, we have $h^1(Z, \mathcal{I}_{E,Z}(x, y)) >$

0. Hence the restriction map $H^0(Z, L|Z) \rightarrow H^0(E, L|E)$ is surjective. Hence if E imposes at most $\deg(E) - 1$ condition to $H^0(W, L)$, then Z imposes at most $\deg(Z) - 1$ condition to $H^0(W, L)$. If $h^i(W, L) = 0$ for all $i \geq 2$, this is equivalent to say that if $h^1(W, \mathcal{I}_E \otimes L) > h^1(W, L)$, then $h^1(W, \mathcal{I}_Z \otimes L) > h^1(W, L)$. We also get $h^1(W, \mathcal{I}_Z \otimes L) \geq h^1(W, \mathcal{I}_E \otimes L)$.

Remark 5. Let W be any projective scheme, Z a closed subscheme of W , D an effective Cartier divisor on W and L a line bundle on W . Let $\text{Res}_D(Z)$ be the closed subscheme of W with $\mathcal{I}_Z : \mathcal{I}_D$ as its ideal sheaf. We have $\text{Res}_D(Z) \subseteq Z$ and an exact sequence

$$0 \rightarrow \mathcal{I}_{\text{Res}_D(Z)} \otimes L(-D) \rightarrow \mathcal{I}_Z \otimes L \rightarrow \mathcal{I}_{Z \cap D, D} \otimes (L|D) \rightarrow 0 \tag{1}$$

From (1) we get

$$h^i(W, \mathcal{I}_Z \otimes L) \leq h^i(W, \mathcal{I}_{\text{Res}_D(Z)} \otimes L(-D)) + h^i(D, \mathcal{I}_{Z \cap D, D} \otimes (L|D))$$

for all $i \geq 0$. If Z is zero-dimensional, then $\deg(Z) = \deg(\text{Res}_D(Z)) + \deg(Z \cap D)$.

Lemma 1. Fix non-negative integers e, x, y such that $x \geq ey$. Fix any $T \in |\mathcal{O}_{X_e}(h)|$ and any $D \in |\mathcal{O}_{X_e}(f)|$. Set $J := T \cup D$. Let $Z \subset J$ be a zero-dimensional scheme such that either $T \cap D \not\subseteq Z_{red}$ or the connected component of Z with $T \cap D$ as its support is reduced.

(i) We have $h^1(X_e, \mathcal{I}_Z(yh + xf)) = h^1(J, \mathcal{I}_{Z, J}(yh + xf))$ and $h^1(X_e, \mathcal{I}_Z(yh + xf)) = \deg(Z) - h^0(X_e, \mathcal{O}_{X_e}(yh + xf)) + h^0(X_e, \mathcal{I}_Z(yh + xf))$.

(ii) $h^1(X_e, \mathcal{I}_Z(yh + xf)) > 0$ if and only if either $\deg(Z \cap D) \geq y + 2$ or $\deg(Z \cap h) \geq x - ey + 2$ or $\deg(Z) = x - (e - 1)y + 2$.

Proof. Part (i) follows from $h^1(X_e, \mathcal{O}_{X_e}(yh + xf)) = h^0(X_e, \mathcal{O}_{X_e}((y - 1)h + (x - 1)f)) = 0$ (Remark 4). We have $h^0(D, \mathcal{O}_D(yh + xf)) = y + 1$, $h^1(D, \mathcal{O}_D((h - 1)y + xf)) = 0$ and $h^0(T, \mathcal{O}_T(yh + xf)) = x - ey + 1$. Since the scheme-theoretic intersection $T \cap D$ is a point, a Mayer-Vietoris exact sequence gives $h^0(J, \mathcal{O}_J(yh + xf)) = x - (e - 1)y + 1$. Hence the “if” part of (ii) follows from Remark 4. Now assume $\deg(Z \cap T) \leq x - ey + 1$, $\deg(Z \cap T) \leq y + 1$. By part (ii) it is sufficient to prove that $h^1(X_e, \mathcal{I}_Z(yh + xf)) = 0$ if $\deg(Z) \neq x - (e - 1)y + 2$. We have $\deg(Z) = \deg(Z \cap T) + \deg(Z \cap D) - \epsilon$ with $\epsilon = 0$ if $T \cap D \in Z_{red}$ and $\epsilon = 1$ if $T \cap D \subset Z$. Hence $\deg(Z) = x - (e - 1)y + 2$ if and only if $\deg(Z \cap D) = y + 1$, $\deg(Z \cap T) = x - ey + 1$ and $D \cap T \not\subseteq Z$. We exclude also this case. First assume $\deg(Z \cap T) = x - ey + 1$. Hence $h^1(T, \mathcal{I}_{Z \cap T}(yh + xf)) = 0$, $i = 0, 1$. Hence (1) gives $h^1(X_e, \mathcal{I}_{\text{Res}_T(Z)}((y - 1)h + xf)) = h^1(X_e, \mathcal{I}_Z(yh +$

xf). Our assumptions imply $\text{Res}_T(Z) \subset D$ and $\deg(\text{Res}_T(Z)) \leq y$. Hence $h^1(D, \mathcal{I}_{\text{Res}_T(Z)}((y-1)h + xf)) = 0$. Since $h^1(X_e, \mathcal{O}_{X_e}((y-1)h + (x-1)f)) = 0$ (Remark 4), we get $h^1(X_e, \mathcal{I}_{\text{Res}_T(Z)}((y-1)h + xf)) = 0$. Remark 5 gives $h^1(X_e, \mathcal{I}_Z(yh + xf)) = 0$. Now assume $\deg(Z \cap T) \leq x - ey$. Since $\deg(D \cap Z) \leq y + 1$, we have $h^1(X_e, \mathcal{I}_{Z \cap D}(yh + xf)) = h^1(D, \mathcal{I}_{Z \cap D}(yh + xf)) = 0$. Remark 5 gives $h^1(X_e, \mathcal{I}_Z(yh + xf)) \leq h^1(X_e, \mathcal{I}_{\text{Res}_D(Z)}(yh + (x-1)))$. We have $\text{Res}_D(Z) = Z \cap D \setminus D \cap T$ and hence $\deg(\text{Res}_D(Z)) \leq x - ey$. Hence the game with D just done gives $h^1(X_e, \mathcal{I}_{\text{Res}_D(Z)}(yh + (x-1))) = 0$. \square

Lemma 2. Fix integers $e > 0$, $x \geq ey$, $z \geq 0$ and a zero-dimensional scheme $Z \subset X_e$ such that $\deg(Z) = z$. Set $c := \deg(Z \cap h)$ and assume that each connected component of Z intersecting h is contained in h . Let η be the set of all $D \in |\mathcal{O}_{X_e}(f)|(\overline{K})$ such that $\deg(D \cap Z)$ such that $t := \deg(D \cap Z)$ is maximal. Set $\eta' := \{D \in \sigma : D \cap h \cap Z = \emptyset\}$. Assume either $z \leq y + 2$ or $z \leq c + \min\{x, y + 2\}$.

(i) If either $c \geq x - ey + 2$ or $t \geq y + 2$ or $\deg((D \cup h) \cap Z) = x - (e - 1)y + 2$, then $h^1(X_e, \mathcal{I}_Z(x, y)) > 0$.

(ii) Assume $c \leq x - ey + 1$ and $t \leq y + 1$. We have $h^1(X_e, \mathcal{I}_Z(yh + xf)) > 0$ if and only if $c = x - ey + 1$, $t = y + 1$ and $\eta' \neq \emptyset$; in this case we have $h^1(X_e, \mathcal{I}_Z(yh + xf)) = 1$, $\eta = \eta'$ has a unique element, D , and $Z \subset D \cup h$. If Z is defined over K , then D is defined over K .

(iii) Assume $c \leq x - ey + 1$, $z - c \leq y + 2$ and either $c \leq x - ey$ or $\eta' = \emptyset$. We have $h^1(X_e, \mathcal{I}_Z(yh + zf)) > 0$ if and only if $t = y + 2$. In this case η has a unique element, D , and $Z \subset h \cup D$. If Z is defined over K , then D is defined over K .

(iv) If $c \geq x - ey + 1$ and $z \leq c + y$, then $h^1(X_e, \mathcal{I}_Z(hy + xf)) = c - x + ey - 1$.

Proof. Since $x \geq ey - 1$, we have $h^1(X_e, \mathcal{O}_{X_e}(yh + xf)) = 0$. Hence $h^1(X_e, \mathcal{I}_Z(yh + xf)) > 0$ if and only if $h^0(X_e, \mathcal{O}_{X_e}(yh + xf)) - \deg(Z)$. For any $D \in |\mathcal{O}_{X_e}(f)|(\overline{K})$ we have $h^0(D, \mathcal{O}_D(yh + xf)) = y + 1$. Hence if $\deg(Z \cap D) \geq y + 2$, then $D \cap Z$ imposes $< \deg(Z \cap D)$ conditions to $H^0(D, \mathcal{O}_{X_e}(yh + xf))$ and hence to $H^0(X_e, \mathcal{O}_{X_e}(yh + xf))$ (Remark 4). Similarly, since $h^0(h, \mathcal{O}_h(yh + xf)) = x - ey + 1$, we get that if $c \geq x - ey + 1$, then $h^1(X_e, \mathcal{I}_Z(hy + xf)) \geq c - x + ey - 1 > 0$ (Remark 4).

(a) Here we assume $\deg(Z \cap h) \leq x - ey + 1$. Fix $D_1 \in |\mathcal{O}_{X_e}(f)|(\overline{K})$ such that $a_1 := \deg(Z \cap D_1)$ is maximal among all $D \in |\mathcal{O}_{X_e}(f)|(\overline{K})$. If $a_1 \geq y + 2$, then we are done. Hence we may assume $a_1 \leq y + 1$. Set $J := h \cup D_1$. Lemma

1 gives $h^1(J, \mathcal{I}_{J \cap Z}(yh + xf)) > 0$ if and only if $T \cap D_1 \not\subseteq Z$, $\deg(D_1 \cap Z) = y + 1$ and $\deg(T \cap Z) = x - ey + 1$. Hence we may assume $h^1(J, \mathcal{I}_{Z \cap J, J}(yh + xf)) = 0$. Set $E := \text{Res}_J(Z)$. Remark 5 gives $h^1(X_e, \mathcal{I}_Z(yh + xf)) \leq h^1(X_e, \mathcal{I}_E((y - 1)h + (x - 1)f))$. By assumption $h \cap E = \emptyset$. Set $Z_1 := E$ and $z_1 := \deg(E)$. Let $D_2 \in |\mathcal{O}_{X_e}(h)|(\overline{K})$ such that $a_2 := \deg(D_1 \cap E)$ is maximal and set $J_2 := T \cup D_2$ and $E_2 := \text{Res}_{D_2}(E)$. We have $\deg(E_2) = z_1 - a_2$ (Remark 5). For all integers $i \geq 3$ define inductively the curves $D_i | \mathcal{O}_{X_e}(h) |(\overline{K})$, the scheme $E_i \subseteq E_{i-1}$ and the integers a_i and z_i in the following way. Fix any $D_i | \mathcal{O}_{X_e}(h) |$ such that $a_i := \deg(D_i \cap E_{i-1})$ is maximal and set $E_i := \text{Res}_{D_i}(E_{i-1})$. The maximality property of the integers a_i gives that the sequence $\{a_i\}$ is non-decreasing. Since $|\mathcal{O}_{X_e}(h)|$ has no base points, we have $E_i = \emptyset$ (and hence $a_j = 0$ for all $j > 0$) if $z_{i-1} \leq 1$. Hence $E_i = \emptyset$ for all $i \gg 0$.

(a1) Here we assume $a_i \leq y - i + 1$ for all $2 \leq i \leq y - 1$. Applying Remark 5 $y - 2$ times we get $h^1(X_e, \mathcal{I}_Z(hy + xf)) \leq h^1(X_e, \mathcal{I}_{E_{y-1}}((x - y)f))$. Since $a_y \leq 1$, we get that $\pi | \mathcal{E}_{y-1}$ is an isomorphism onto a degree $\deg(E_{y-1})$ subscheme of \mathbb{P}^1 . Since $\pi^*(H^0(\mathbb{P}^1, \mathcal{O}_{\mathbb{P}^1}(x - y))) \rightarrow H^0(X_e, \mathcal{O}_{X_e}((x - y)f))$ is an isomorphism, we get $h^1(X_e, \mathcal{I}_{E_{y-1}}((x - y)f)) = 0$ if $\deg(E_{y-1}) \leq x - y + 1$. Hence $h^1(X_e, \mathcal{I}_Z(hy + xf)) = 0$ if $z \leq c + x$ (under the assumption $a_i \leq y - i + 1$ for all $2 \leq i \leq y$).

(a2) Assume the existence of an integer $i \in \{2, \dots, y - 1\}$ such that $a_i \geq y - i + 2$ and call w the minimal such an integer. First assume $w \leq e$. Set $A' := D_2 + \dots + D_{w-1}$ and $A := A' \cup J$ (as effective divisor of X_e , not as a set, because we may have $D_i = D_j$ even if $i \neq j$). We have $h^1(X_e, \mathcal{I}_{A' \cap Z}(((y - 1)h + (x - 1)f))) = 0$ and hence $h^1(X_e, \mathcal{I}_{A \cap Z}(yh + xf)) = 0$ (Remark 5). Thus $h^1(X_e, \mathcal{I}_Z(yh + xf)) \leq h^1(X_e, \mathcal{I}_{\text{Res}_A(Z)}(((y - 1)h + (x - w)f)))$. We have $\deg(A \cap Z) = c + a_1 + \dots + a_w \geq c + wa_w \geq c + w(y - w + 4)$. The function $\phi : [0, y + 4] \rightarrow \mathbb{R}$ defined by $\phi(t) = t(y + 4 - t)$ is increasing for $0 \leq t \leq (y + 4)/2$ and decreasing for $(y + 4)/2 \leq t \leq y + 4$. First assume $w = 3$, i.e. $a_1 = a_2 = y + 1$. Hence $z \geq c + 2y + 2$, absurd. Now assume $w = y - 1$. We get $\deg(Z) \geq c + 5y - 5$, absurd. Now assume $w > e$. We get $z \geq c + w(y + 4 - w)$. Since $w \leq y - 1$ and either $z \leq y + 2$ or $z \leq c + y$, we get a contradiction (even in the case $w = y - 1$).

(a3) Since \overline{K} is a Galois extension of K , the uniqueness of D in (ii) and (iii) implies that D is defined over K .

(b) Assume $c \geq x - ey + 1$ and $z \leq c + y$. We have $\deg(\text{Res}_h(Z)) = z - c$. Our assumptions on Z gives $\text{Res}_h(Z) \cap h = \emptyset$. By step == we get $h^1(X_e, \mathcal{I}_{\text{Res}_h(Z)}((y - 1)h + xf)) = 0$. Hence (1) gives $h^1(X_e, \mathcal{I}_Z(hy + xf)) = c - x + ey - 1$. □

Proof of Theorem 1. By the definition of \mathcal{C}^\perp the minimum distance of \mathcal{C}^\perp is the minimal cardinality of a set $S \subseteq B$ such that $h^0(C, \mathcal{O}_C(yh+xf)(-E-S)) > h^0(C, \mathcal{O}_C(yh+xf)(-E)) - \#(S)$, i.e. such that $h^1(C, \mathcal{O}_C(yh+xf)(-E-S)) > h^1(C, \mathcal{O}_C(yh+xf)(-E))$ (Riemann-Roch) Since $m \leq x - ey + 1$ and $s \leq y + 1$, Lemma 2 gives $h^1(C, \mathcal{O}_C(yh+xf)(-E)) = h^1(C, \mathcal{O}_C(yh+xf))$. Hence the minimum distance of \mathcal{C}^\perp is the minimal cardinality of a set $S \subseteq B$ such that $h^1(C, \mathcal{O}_C(yh+xf)(-E-S)) > h^1(C, \mathcal{O}_C(yh+xf))$. By Remarks 2 and 3 this is the minimal cardinality of a subset $S \subseteq B$ such that

Fix any $S \subseteq B$ supporting a codeword of \mathcal{C}^\perp with minimum codeword. There is a unique, up to a non-zero scalar, codeword with S as its support. \square

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