

NEW CONSTRUCTIONS OF FUNCTION AND SET CHAINS

Martin Dowd

1613 Wintergreen Pl.
Costa Mesa, CA 92626, USA

Abstract: Function and set chains can be constructed using ordinal notation systems. Here, earlier results are improved, and longer chains constructed. In particular, a simpler method is given for constructing set chains.

AMS Subject Classification: 03E55

Key Words: Galvin-Hajnal order, stationary reflection order

1. Function Chains for Schemes

Function chains for schemes have been given in [4]. These will be reviewed here, with some changes. Recall from [6] that a scheme over κ where $\kappa \in \text{Inac}$ is a pair $\sigma = \langle \sigma, \phi \rangle$ where $\sigma < \kappa^+$ and ϕ is a function whose domain is the set of limit ordinals $\alpha \leq \sigma$. For $\alpha \in \text{Dom}(\phi)$, $\phi(\alpha)$ is an increasing function with domain an ordinal $\eta \leq \kappa$, and whose range is an unbounded subset of α . If $\text{Cf}(\alpha) < \kappa$ then $\eta < \kappa$, and if $\text{Cf}(\alpha) = \kappa$ then $\eta = \kappa$.

Recursions and inductions on schemes are broken into four cases, as follows.

Case 0: $\sigma = 0$.

Case 1: $\sigma = \tau + 1$. τ denotes $\sigma_{\leq \tau}$.

Case 2: $\sigma \in \text{Lim}$, $\text{Cf}(\sigma) < \kappa$. For $\xi < \eta$ σ_ξ denotes $\phi(\sigma)(\xi)$; and σ_ξ denotes $\sigma_{\leq \sigma_\xi}$.

Case 3: $\sigma \in \text{Lim}$, $\text{Cf}(\sigma) = \kappa$. σ_ξ and σ_ξ are as in case 2, for $\xi < \kappa$.

The thin subset $T_\sigma \subseteq \kappa$ is defined recursively as follows.

- 0: \emptyset .
- 1: T_τ .
- 2: $(\eta + 1) \cup \cup_{\xi < \eta} T_{\sigma_\xi}$.
- 3: $\nabla_{\xi < \kappa} T_{\sigma_\xi}$.

A function $f_\sigma : \text{In}_\kappa \mapsto \kappa$ is defined recursively as follows.

- 0: The identically 0 function.
- 1: $f_\tau + 1$.
- 2: $\sup_{\xi < \eta} f_{\sigma_\xi}$, with $f_\sigma(\lambda)$ set to 0 if $\lambda \leq \theta$.
- 3: $\text{dsup}_{\xi < \kappa} f_{\sigma_\xi}$.

This definition was given in [4], for domain κ , and it was observed that these functions are the well-known canonical functions, which are unique mod the thin ideal. As noted in [7], the domain In_κ is convenient for defining f_{κ^+} . This results in the triviality that when λ is the smallest inaccessible, f_σ is the empty function for all σ . The proviso in case 2 ensures that $f_\sigma(\lambda) < \lambda^+$, as a simple induction shows; this does not seem to be necessary, though.

Lemma 1. *Suppose σ', σ are schemes, and $\lambda \notin T_{\sigma'} \cup T_\sigma$.*

- a. *If $\sigma' \leq \sigma$ then $f_{\sigma'}(\lambda) \leq f_\sigma(\lambda)$.*
- b. *If $\sigma' = \sigma$ then $f_{\sigma'}(\lambda) = f_\sigma(\lambda)$.*
- c. *If $\sigma = \tau + 1$ then $f_\sigma(\lambda) = f_\tau(\lambda) + 1$.*
- d. *If $\sigma' < \sigma$ then $f_{\sigma'}(\lambda) < f_\sigma(\lambda)$.*

Proof. Part a follows by induction on σ . In case 0, σ is 0 also and the claim is trivial. In case 1, if $\sigma' \leq \tau$ then inductively $f_{\sigma'}(\lambda) \leq f_\tau(\lambda)$, and clearly $f_\tau(\lambda) < f_\sigma(\lambda)$. If $\sigma' = \sigma$ then inductively $f_{\sigma'}(\lambda) = f_\tau(\lambda)$, and $f_{\sigma'}(\lambda) = f_\sigma(\lambda)$ follows. In case 2, if $\sigma' < \sigma$ then $\sigma' < \sigma_\xi$ for some $\xi < \eta$. Inductively, $f_{\sigma'}(\lambda) \leq f_{\sigma_\xi}(\lambda)$, and clearly $f_{\sigma_\xi}(\lambda) \leq f_\sigma(\lambda)$. If $\sigma' = \sigma$, by the argument just given, $f_{\sigma'_\xi}(\lambda) \leq f_\sigma(\lambda)$ for all $\xi < \eta'$, whence $f_{\sigma'}(\lambda) \leq f_\sigma(\lambda)$. The argument for case 3 is similar to that for case 2. Part b is immediate from part a. For part c let $\tau_1 = \sigma_{\leq \tau}$ and use part b. Part d follows by parts a and c. □

A scheme $\sigma \downarrow \lambda$ over λ , for $\lambda \notin T_\sigma$, may be defined recursively as follows. Write $\sigma \downarrow \lambda$ as $\langle \sigma', \phi' \rangle$.

- 0: The scheme with $\sigma' = 0$.
- 1: $\tau \downarrow \lambda$ with τ' replaced by $\tau' + 1$.
- 2: $\sqcup_{\xi < \eta} \sigma_\xi \downarrow \lambda$, with $\phi'(\sigma')(\xi)$ set to $(\phi(\sigma)(\xi))'$.
- 3: $\sqcup_{\xi < \lambda} \sigma_\xi \downarrow \lambda$, with $\phi'(\sigma')(\xi)$ set to $(\phi(\sigma)(\xi))'$.

Lemma 2. *Suppose σ is a scheme, $\lambda \in \text{In}_\kappa$, and $\lambda \notin T_\sigma$.*

- a, $f_\sigma \upharpoonright \lambda = f_{\sigma \downarrow \lambda}$.

b, $f_\sigma(\lambda) = \sigma \downarrow \lambda$.

Proof. This is lemma 5 of [5]; for convenience the details are given. For part a, the cases of the induction are as follows, where $\mu \in \text{In}_\lambda$.

- 0: $f_0(\mu) = 0 = f_{0 \downarrow \lambda}(\mu)$.
- 1: $f_\sigma(\mu) = f_\tau(\mu) + 1 = f_{\tau \downarrow \lambda}(\mu) + 1 = f_{\sigma \downarrow \lambda}(\mu)$.
- 2: $f_\sigma(\mu) = \sup_{\xi < \eta} f_{\sigma_\xi}(\mu) = \sup_{\xi < \eta} f_{\sigma_\xi \downarrow \lambda}(\mu) = f_{\sigma \downarrow \lambda}(\mu)$.
- 3: $f_\sigma(\mu) = \text{dsup}_{\xi < \kappa} f_{\sigma_\xi}(\mu) = \text{dsup}_{\xi < \lambda} f_{\sigma_\xi \downarrow \lambda}(\mu) = f_{\sigma \downarrow \lambda}(\mu)$.

For part b, the cases of the induction are as follows.

- 0: $f_0(\lambda) = 0 = 0 \downarrow \lambda$.
- 1: $f_\sigma(\lambda) = f_\tau(\lambda) + 1 = \tau \downarrow \lambda + 1 = \sigma \downarrow \lambda$.
- 2: $f_\sigma(\lambda) = \sup_{\xi < \eta} f_{\sigma_\xi}(\lambda) = \sup_{\xi < \eta} \sigma_\xi \downarrow \lambda = \sigma \downarrow \lambda$.
- 3: $f_\sigma(\lambda) = \text{dsup}_{\xi < \kappa} f_{\sigma_\xi}(\lambda) = \text{dsup}_{\xi < \lambda} \sigma_\xi \downarrow \lambda = \sigma \downarrow \lambda$.

□

2. Infinitary Veblen Functions

The infinitary Veblen functions were defined in [11]. Further work on them includes [9] and [10]. A certain set of symbols for them are called Klammer-symbols, or Schutte brackets. For convenience a self-contained treatment will be given here, following [11].

For basic facts about normal functions $f : \text{Ord} \mapsto \text{Ord}$, and club subclasses $R \subseteq \text{Ord}$, see [8], section 4.2. In particular, the following hold.

- If f is a normal function then $\text{Ran}(f)$ is a club subclass.
- If R is a club subclass then the function $\text{Enum}(R)$ enumerating R in natural order is normal.
- If f is normal then the function $\text{Fix}(f)$ enumerating the subclass $\{\alpha : f(\alpha) = \alpha\}$ of fixed points of f is normal.
- The intersection of a set of club class is club.
- Given a normal function f , suppose $f(\alpha) > \alpha$. Let $\alpha_0 = \alpha$, and for $n \geq 0$ let $\alpha_{n+1} = f(\alpha_n)$. Then $\sup_n \alpha_n$ is the smallest fixed point of f which is greater than α .
- If $f(0) > 0$ then $\text{Fix}(f)(0) > f(0)$.

Let \mathcal{A} be the set of ordinal valued sequences $\alpha_0, \dots, \alpha_\mu$ for some μ , where only finitely many α_ξ are nonzero, and $\alpha_\mu > 0$ if $\mu > 0$. \mathcal{A} may be ordered in reverse lexicographic order, where if $\bar{\alpha} = \alpha_0, \dots, \alpha_\mu$ and $\bar{\alpha}' = \alpha'_0, \dots, \alpha'_{\mu'}$, then $\bar{\alpha}' <_{\text{rl}} \bar{\alpha}$ iff $\mu' < \mu$, or $\mu' = \mu$, and for some $\gamma \leq \mu$, $\alpha'_\gamma < \alpha_\gamma$ and $\alpha'_\xi = \alpha_\xi$ for $\gamma < \xi \leq \mu$. This relation is well-known to be a well-order (see [11] for example).

Let \mathcal{A}_* be the sequences obtained from the sequences of \mathcal{A} , by replacing some α_γ by $*$, provided $\alpha_\delta = 0$ for $\delta < \gamma$. The order $<_{r1}$ may be extended to $\mathcal{A} \cup \mathcal{A}_*$, by taking $\alpha < *$ for $\alpha \in \text{Ord}$.

Let θ be a cardinal. A function $\phi : \mathcal{A} \mapsto \text{Ord}$ will be defined, together with functions $\phi_{\tilde{\alpha}} : \text{Ord} \mapsto \text{Ord}$ for each $\tilde{\alpha} \in \mathcal{A}_*$. The cases of the recursion are as follows.

1. $\phi_{\tilde{\alpha}}(\zeta) = \phi(S_1(\tilde{\alpha}, \zeta))$, where in $S_1(\tilde{\alpha}, \zeta)$, $*$ is replaced by ζ (and trailing 0's removed if $\tilde{\alpha}_\mu = *$ and $\zeta = 0$).
2. For $\alpha \in \text{Ord}$, $\phi(\alpha) = \theta^\alpha$.
3. For $\tilde{\alpha} \in \mathcal{A}_*$ with $\mu > 0$, $\tilde{\alpha}_0 = *$, $\tilde{\alpha}_\gamma = 0$ for $0 < \gamma < \nu$, and $\tilde{\alpha}_\nu > 0$, $\phi_{\tilde{\alpha}} = \text{Enum}(\cap_{\beta < \alpha_\nu, \gamma < \nu} \text{Ran}(\text{Fix}(\phi_{S_2(\alpha, \beta, \gamma)})))$, where in $S_2(\alpha, \beta, \gamma)$, α_0 is replaced by 0 and α_γ by $*$ (this is a null operation if $\gamma = 0$), and α_ν is replaced by β (and trailing 0's removed if $\nu = \mu$ and $\beta = 0$).

Lemma 3. For $\tilde{\alpha} \in \mathcal{A}_*$, $\phi_{\tilde{\alpha}}$ is normal.

Proof. The proof is by induction on $<_{r1}$. For the basis, $\tilde{\alpha} = *$; the lemma follows since $\zeta \mapsto \theta^\zeta$ is normal. For the induction step, if $\tilde{\alpha}_0 = *$ the lemma follows by the induction hypothesis and the definition of ϕ . Otherwise, using obvious notation write $\tilde{\alpha}$ as $0\bar{0}*\bar{\eta}$; then $\phi_{0\bar{0}*\bar{\eta}}(\zeta) = \phi_{*\bar{0}\zeta\bar{\eta}}(0)$. Suppose $\xi < \zeta$. Directly from the definition of ϕ , $\text{Ran}(\phi_{*\bar{0}\xi\bar{\eta}}) \supseteq \text{Ran}(\phi_{*\bar{0}\zeta\bar{\eta}})$ and $\phi_{*\bar{0}\xi\bar{\eta}}(0) < \phi_{*\bar{0}\zeta\bar{\eta}}(0)$. In particular, $\phi_{0\bar{0}*\bar{\eta}}(\xi) < \phi_{0\bar{0}*\bar{\eta}}(\zeta)$; that is, $\phi_{\tilde{\alpha}}$ is increasing.

For $\zeta \in \text{Lim}$ let $\chi = \sup_{\xi < \zeta} \phi_{\tilde{\alpha}}(\xi)$. By what was just proved, $\chi \leq \phi_{\tilde{\alpha}}(\zeta)$, whence $\chi \leq \phi_{*\bar{0}\zeta\bar{\eta}}(0)$. Thus, to show that $\phi_{\tilde{\alpha}}$ is continuous it suffices to show that $\chi \in \text{Ran}(\phi_{*\bar{0}\zeta\bar{\eta}})$, since then $\chi = \phi_{*\bar{0}\zeta\bar{\eta}}(0) = \phi_{\tilde{\alpha}}(\zeta)$. For this in turn it suffices to show that $\phi_{S_2(*\bar{0}\zeta\bar{\eta}, \beta, \gamma)}(\chi) = \chi$ for $\beta < \zeta$, $\gamma < \nu$.

Suppose $\beta < \xi$, and for ease of notation let

$$\tilde{\alpha}_1 \text{ denote } S_2(*\bar{0}\zeta\bar{\eta}, \beta, \gamma), \text{ and } \tilde{\alpha}_2 \text{ denote } *\bar{0}\zeta\bar{\eta}.$$

From the definition of ϕ , $\text{Ran}(\phi_{\tilde{\alpha}_2}) \subseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_1}))$, and so $\phi_{\tilde{\alpha}_2}(0) \in \text{Fix}(\phi_{\tilde{\alpha}_1})$, whence $\phi_{\tilde{\alpha}_1}(\phi_{\tilde{\alpha}_2}(0)) = \phi_{\tilde{\alpha}_2}(0)$. Using this fact,

$$\begin{aligned} \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\chi) &= \\ \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\sup_{\xi < \zeta} \phi_{\tilde{\alpha}}(\xi)) &= \sup_{\xi < \zeta} \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\phi_{\tilde{\alpha}}(\xi)) = \\ \sup_{\beta < \xi < \zeta} \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\phi_{\tilde{\alpha}}(\xi)) &= \sup_{\beta < \xi < \zeta} \phi_{\tilde{\alpha}}(\xi) = \chi. \end{aligned} \quad \square$$

Lemma 4. If $\tilde{\beta} <_{r1} \tilde{\alpha}$ then $\text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \cup \{1\} \supseteq \text{Ran}(\phi_{\tilde{\alpha}})$.

Proof. Writing $\tilde{\alpha}$ as $0\bar{0}*\bar{\alpha}$, $\phi_{0\bar{0}*\bar{\alpha}}(\zeta) = \phi_{*\bar{0}\zeta\bar{\alpha}}(0)$, provided $\zeta \neq 0$ if $\bar{\alpha}$ is null. By the definition of ϕ and induction on $<_{r1}$, $\phi_{\tilde{\alpha}}(\zeta) \in \text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \cup \{1\}$, proving the lemma. □

Let $\mathcal{A}_*^0 = \{\tilde{\alpha} \in \mathcal{A}_* : \tilde{\alpha}_0 = *\}$. The following are readily seen, for $\tilde{\beta} <_{r1} \tilde{\alpha}$.

- For $\tilde{\alpha} \in \mathcal{A}_*^0$, $1 \in \text{Ran}(\phi_{\tilde{\alpha}})$ iff $\tilde{\alpha} = *$.
- For $\tilde{\alpha}, \tilde{\beta} \in \mathcal{A}_*^0$, If $\tilde{\beta} <_{r1} \tilde{\alpha}$ then $\text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \supseteq \text{Ran}(\phi_{\tilde{\alpha}})$.
- For $\tilde{\alpha}, \tilde{\beta} \in \mathcal{A}_*$, $\text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \supseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}}))$.

For $\mu \in \text{Ord}$ let $\tilde{\alpha}_\mu$ be the element of \mathcal{A}_* where $\tilde{\alpha}_{\mu 0} = *$, $\tilde{\alpha}_{\mu\mu} = 1$, and $\tilde{\alpha}_{\mu\xi} = 0$ for $0 < \xi < \mu$. Let ψ be the function $\mu \mapsto \tilde{\alpha}_\mu(0)$.

Lemma 5. *The function ψ is normal.*

Proof. For $\mu \in \text{Ord}$ let $\tilde{\alpha}_\mu^*$ be the element of \mathcal{A}_* where $\tilde{\alpha}_{\mu\mu} = *$ and $\tilde{\alpha}_{\mu\xi} = 0$ for $0 \leq \xi < \mu$. From the definition of ϕ , $\phi_{\tilde{\alpha}_\mu^*} = \text{Enum}(\cap_{\gamma < \mu} \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_\gamma^*}))$). It follows from this and lemma 4 that $\mu \mapsto \tilde{\alpha}_\mu(0)$ is increasing. For $\mu \in \text{Lim}$ let $\chi = \sup_{\gamma < \mu} \phi_{\tilde{\alpha}_\gamma}(0)$. From the definition of ϕ , for $\gamma < \delta$ $\text{Ran}(\phi_{\tilde{\alpha}_\delta}) \subseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_\gamma^*}))$. Using this fact, it follows by arguments as given above that $\phi_{\tilde{\alpha}_\mu}(0) = \chi$. \square

The least fixed point of ψ is called the “large Veblen ordinal”. Notation varies; Λ_0 will be used here. Clearly, $\Lambda_0 = \phi_{\tilde{\alpha}_{\Lambda_0}}(0) = \sup_{\mu < \Lambda_0} \phi_{\tilde{\alpha}_\mu}(0)$. It follows that for $\mu < \Lambda_0$, Λ_0 is a fixed point of $\phi_{\tilde{\alpha}_\mu}$. Using lemma 4, if $\tilde{\alpha} < \tilde{\alpha}_{\Lambda_0}$ then Λ_0 is a fixed point of $\phi_{\tilde{\alpha}}$, whence if $\zeta < \Lambda_0$ then $\phi_{\tilde{\alpha}}(\zeta) < \Lambda_0$.

Say that $\bar{\alpha} \in \mathcal{A}$ is maximal if $\phi(\tilde{\beta}) \neq \phi(\bar{\alpha})$ whenever $\beta >_{r1} \alpha$.

Lemma 6. *For $\bar{\alpha} \in \mathcal{A}$, let $\nu_1 > \dots > \nu_k$ be the indices such that $\alpha_{\nu_i} > 0$.*

- a. *For $i > 1$, $\phi(\bar{\alpha}) > \alpha_{\nu_i}$.*
- b. *$\phi(\bar{\alpha}) > \alpha_{\nu_1}$ iff $\bar{\alpha}$ is maximal.*
- c. *If $\phi(\bar{\alpha}) \notin \text{Ran}(\psi)$ then $\nu_k < \phi(\bar{\alpha})$.*
- d. *If $\tau \in \text{Ran}(\phi_*)$ then $\tau = \phi(\bar{\alpha})$ for a unique maximal $\bar{\alpha}$.*

Proof. For $1 \leq j < i$ let $\bar{\alpha}_i$ be $\bar{\alpha}$ with α_{ν_j} replaced by 0. It is easily seen that $\phi(\bar{\alpha}_{i-1}) > \phi(\bar{\alpha}_i) \geq \alpha_{\nu_i}$ for $i > 1$, and part a follows. Write $\bar{\alpha}$ as $\bar{0}\alpha_\nu\alpha_{\nu+1}\bar{\alpha}'$, where for ease of notation $\mu > \nu$ is assumed. Let $\tilde{\alpha} = \bar{0}*\alpha_{\nu+1}\bar{\alpha}'$. Since $\phi_{\tilde{\alpha}}$ is normal, $\phi(\alpha) \geq \alpha_\nu$. Let $\tilde{\beta} = *\bar{0}0, \alpha_{\nu+1} + 1, \bar{\alpha}'$. Then $\text{Ran}(\phi_{\tilde{\beta}}) = \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}}))$, and so $\phi(\bar{\alpha}) = \alpha_\nu$ iff $\alpha_\nu \in \text{Ran}(\phi_{\tilde{\beta}})$ iff $\bar{\alpha}$ is not maximal. The case $\nu = \mu$ is similar, proving part b. For part c, let $\gamma \in \text{Ran}(\psi)$ be least such that $\phi(\bar{\alpha}) < \gamma$. Then $\phi(\bar{\alpha}_{\nu_k}) \leq \phi(\bar{\alpha}) < \gamma$, whence $\nu_k < \phi(\bar{\alpha}_{\nu_k}) \leq \phi(\bar{\alpha})$. For part d, there is a μ such that $\tau < \phi_{\tilde{\alpha}_\mu}(0) \leq \phi_{\tilde{\alpha}_\mu}(\tau)$, and so there is a lexicographically least $\tilde{\alpha}$ such that $\tau < \phi_{\tilde{\alpha}}(\tau)$.

If $\tilde{\alpha} = *$ then $\tau = \phi(\zeta)$ for some ζ by hypothesis, and $\zeta < \tau$ since $\zeta = \tau$ is impossible. Now, $\tau = \phi_{\tilde{\beta}}(\tau)$ for $\tilde{\beta} <_{r1} \tilde{\alpha}$. If $\tilde{\alpha}_0 = *$ then by the definition of ϕ $\tau = \phi_{\tilde{\alpha}}(\zeta)$ for some ζ , and $\zeta = \tau$ is impossible. Otherwise $\tilde{\alpha} = 0\bar{0} * \bar{\alpha}'$, and $\tau < \phi_{0\bar{0}*\bar{\alpha}'}(\tau) = \phi_{*\bar{0}\tau\bar{\alpha}'}(0)$; but then by the definition of ϕ $\tau = \phi_{*\bar{0}\tau\bar{\alpha}'}(\zeta)$ for some ζ , which is impossible. \square

An infinitary Veblen function is a function $\phi_k(\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k)$ where $k \geq 1$. If $\xi_1 < \dots < \xi_k$ and $\alpha_i > 0$ the arguments are said to be proper; the value of ϕ_k is $\phi(\bar{\alpha})$ where $\bar{\alpha}$ is the sequence $\zeta, \dots, \alpha_1 \dots \alpha_2 \dots \alpha_k$, where α_i occurs in position $1 + \zeta_i$, and all other elements of $\bar{\alpha}$ are 0. If the arguments are improper the value is 0.

Corollary 7. *For $\tau \in \text{Ran}(\phi_*)$ with $\tau < \Lambda_0$ there are a unique k and values $\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k < \tau$ such that $\tau = \phi_k(\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k)$.*

Proof. Immediate. □

Let O denote the function with ordinal arguments, where $O(\alpha_1, \alpha_2)$ equals -1 if $\alpha_1 < \alpha_2$, 0 if $\alpha_1 = \alpha_2$, and $+1$ if $\alpha_1 > \alpha_2$,

Lemma 8. *Suppose $\tilde{\alpha}_1, \tilde{\alpha}_2 \in \mathcal{A}_*^0$ and $\tilde{\alpha}_1 <_{rl} \tilde{\alpha}_2$. Then $O(\phi_{\tilde{\alpha}_1}(\zeta_1), \phi_{\tilde{\alpha}_2}(\zeta_2)) = O(\zeta_1, \phi_{\tilde{\alpha}_2}(\zeta_2))$.*

Proof. $\text{Ran}(\phi_{\tilde{\alpha}_2}) \subseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_1}))$, so $\phi_{\tilde{\alpha}_1}(\phi_{\tilde{\alpha}_2}(\zeta_2)) = \phi_{\tilde{\alpha}_2}(\zeta_2)$; the lemma follows. □

3. IV_o Terms

As in [7], a CNF function is a function $C_k(\eta_k, \sigma_k, \dots, \eta_1, \sigma_1)$, where $k \geq 1$. If $\eta_k > \dots > \eta_1$ and $0 < \sigma_i < \theta$ the arguments are said to be proper, and the value of C_k is $\theta^{\eta_k} \cdot \sigma_k + \dots + \theta^{\eta_1} \cdot \sigma_1$. If the arguments are improper, the value is defined to be 0.

Let IV_o be the system of terms, whose leaves are ordinals σ , and whose interior nodes are either a CNF function C_k or an infinitary Veblen function ϕ_k . A proper term is one where the arguments are proper at each interior node. A normal form term is a proper term where at each interior node, the value of the function is greater than the value of any of its arguments.

Theorem 9. *If $\alpha < \Lambda_0$ then there is a unique normal form IV_o term whose value is α .*

Proof. Let $\kappa^{+\eta_k} \cdot \sigma_k + \dots + \kappa^{+\eta_1} \cdot \sigma_1$ be the Cantor normal form for α . The proof is by induction on α , with cases as follows, where t denotes the term.

Case 0: $\alpha = 0$. Then $t = 0$.

Case 1: $k = 1$ and $\eta_1 = 0$. Then $t = \sigma_1$.

Case 2: $k > 1$ or $\sigma_1 > 1$ or $\eta_1 > 0$ and $\eta_1 < \kappa^{+\eta}$. Then $t = C_k(t_{\eta_k}, t_{\sigma_k}, \dots, t_{\eta_1}, t_{\sigma_1})$.

Case 3: $k = 1$ and $\sigma_1 = 1$ and $\eta_1 = \kappa^{+\eta}$. In this case, $t = \phi_k(t_\zeta, t_{\xi_1}, t_{\alpha_1}, \dots, t_{\xi_k}, t_{\alpha_k})$ where $\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k < \alpha$ are the unique values as in lemma 7. \square

If t_i is a term with value α_i for $i = 1, 2$, write $O(t_1, t_2)$ for $O(\alpha_1, \alpha_2)$. Let $P(t)$ equal 1 if t is proper, else 0 Let $N(t)$ equal 1 if t is normal, else 0.

- Lemma 10.**
- a. $O(t_1, t_2)$ depends only on the values $O(\sigma_1, \sigma_2)$ and $P(\sigma)$ for the leaves of t_1, t_2 .
 - b. $P(t)$ depends only on the values $O(\sigma_1, \sigma_2)$ and $P(\sigma)$ for the leaves of t .
 - c. $N(t)$ depends only on the values $O(\sigma_1, \sigma_2)$ and $P(\sigma)$ for the leaves of t .

Proof. Parts a and b are proved by induction on the total number of nodes. For part b, if t is a leaf then $P(t) = 1$. If t is a CNF function, $P(t)$ iff $P(s)$ holds for all subterms s , and certain inequalities hold among the subterms. If t is a IV function the argument is similar.

Part a falls into the following cases. If t_1 a leaf and t_2 a leaf then $O(t_1, t_2)$ is known by hypothesis. Suppose t_1 an ordinal and t_2 is a CNF function. If t_2 is improper then $O(t_1, t_2)$ equals 0 if $t_1 = 0$, else 1. If t_2 is proper then $O(t_1, t_2)$ equals -1 , unless $t_2 = \kappa^{+0} \cdot \sigma_2$, in which case it equals $O(t_1, \sigma_2)$. Suppose t_1 an ordinal and t_2 is an IV function. If t_2 is improper then $O(t_1, t_2)$ equals 0 if $t_1 = 0$, else 1. If t_2 is proper then $O(t_1, t_2) = -1$. In the remaining cases, if t_1 and t_2 are both improper then $O(t_1, t_2) = 0$, if t_1 is improper and t_2 is proper then $O(t_1, t_2) = -1$, and if t_1 is proper and t_2 is improper then $O(t_1, t_2) = 1$; so it may be assumed that t_1 and t_2 are proper. If t_1 and t_2 are both CNF functions, $O(t_1, t_2)$ is determined by the lexicographic order, using the induction hypothesis. Suppose t_1 is a CNF function and t_2 is an IV functions. If $\eta_k < \alpha_2$ (where η_k is the leading exponent in t_1) then $\alpha_1 < \alpha_2$. If $\eta_k > \alpha_2$ then $\alpha_1 > \alpha_2$. If $\eta_k = \alpha_2$ then $\alpha_1 > \alpha_2$, unless $k = 1$ and $\sigma_k = 1$, in which case $\alpha_1 = \alpha_2$. If t_1 is an IV function and t_2 is a CNF functions, reverse the roles of t_1 and t_2 in the preceding argument. Suppose t_1 and t_2 are IV functions, say $\phi_{\tilde{\alpha}_1}(\zeta_1)$ and $\phi_{\tilde{\alpha}_2}(\zeta_2)$, If $\tilde{\alpha}_1 <_{r1} \tilde{\alpha}_2$ (which may be determined from the arguments using the induction hypothesis) then $O(t_1, t_2) = O(\zeta_1, \alpha_2)$. If $\tilde{\alpha}_1 = \tilde{\alpha}_2$ then $O(t_1, t_2) = O(\zeta_1, \zeta_2)$. If $\tilde{\alpha}_2 <_{r1} \tilde{\alpha}_1$ then $O(t_1, t_2) = O(\alpha_1, \zeta_2)$.

Part c now follows inductively. The claim is trivial if t is a leaf. If t is a CNF function then t is normal if its subterms are normal, except if $k = 1$ and $\eta_k = 0$. If t is an IV function then t is normal if its subterms are normal, except if $t = \phi_{\tilde{\alpha}_1}(\phi_{\tilde{\alpha}_2}(\zeta_2))$ where $\tilde{\alpha}_1 <_{r1} \tilde{\alpha}_2$. \square

4. IV Scheme Terms

From hereon the cardinal θ will be κ^+ for some $\kappa \in \text{Inac}$. A scheme term over κ is defined to be a term, whose leaves are 0 or a scheme over κ with $\sigma > 0$; and whose interior nodes are either a CNF function C_k or an infinitary Veblen function ϕ_k where $k \geq 1$. The value α of a term α is the value of the IV_o term obtained by replacing each scheme σ by its value σ . Clearly α depends only on the values of the schemes.

It follows from theorem 9 that for each $\alpha < \Lambda_0$ there is a normal form term α , whose ordinal is α . α is unique, up to the choice of schemes σ for the leaf ordinals σ . From hereon, unless specifically stated otherwise, IV scheme terms will be assumed to be in normal form.

For a term α let T_α be the union over the σ occurring at leaves of α of the T_σ .

Let $\alpha \downarrow \lambda$ be α , with each leaf σ replaced by $\sigma \downarrow \lambda$, each node C_k replaced by $C_{\lambda k}$, and each node ϕ_k replaced by $\phi_{\lambda k}$. Given α , the notation $\alpha \downarrow \lambda$ will be used to denote the ordinal specified by $\alpha \downarrow \lambda$.

Lemma 11. *Suppose α is a scheme, $\lambda \in \text{In}_\kappa$, and $\lambda \notin T_\alpha$.*

- a, $f_\alpha \upharpoonright \lambda = f_{\alpha \downarrow \lambda}$.
- b, $f_\alpha(\lambda) = \alpha \downarrow \lambda$.

Proof. For schemes this is lemma 2. Let the argument list of a function be denoted A_1, \dots, A_l . For part a, suppose $\mu < \lambda$. For CNF functions, $f_{\alpha \downarrow \lambda}(\mu) = f_{C_k(A_1 \downarrow \lambda, \dots, A_l \downarrow \lambda)}(\mu) = C_{\mu k}(f_{A_1 \downarrow \lambda}(\mu), \dots, f_{A_l \downarrow \lambda}(\mu)) = C_{\mu k}(f_{A_1}(\mu), \dots, f_{A_l}(\mu)) = f_{C_k(A_1, \dots, A_l)}(\mu) = f_\alpha(\mu)$. The proof for IV functions is similar. For part b, for CNF functions, $f_\alpha(\lambda) = f_{C_k(A_1, \dots, A_l)}(\lambda) = C_{\lambda k}(f_{A_1}(\lambda), \dots, f_{A_l}(\lambda)) = C_{\lambda k}(A_1 \downarrow \lambda, \dots, A_l \downarrow \lambda) = C_k(A_1, \dots, A_l) \downarrow \lambda = \alpha \downarrow \lambda$. The proof for IV functions is similar. □

5. Function Chains for IV Terms

For each IV term α a function $f_\alpha : \text{In}_\kappa \mapsto \kappa$ may be defined by recursion on α . The cases of the recursion are numbered as in theorem 9, and are as follows. In cases 2 and 3, $C_{\lambda k}$ and $\phi_{\lambda k}$ is written for the function over the cardinal λ^+ .

- 0. $f_0 = 0$
- 1. f_σ where α is the scheme σ .
- 2. $f_\alpha(\lambda) = C_{\lambda k}(f_{\eta_k}(\lambda), f_{\sigma_k}(\lambda), \dots, f_{\eta_1}(\lambda), f_{\sigma_1}(\lambda))$.
- 3. $f_\alpha(\lambda) = \phi_{\lambda k}(f_\zeta(\lambda), f_{\xi_1}(\lambda), f_{\alpha_1}(\lambda), \dots, f_{\xi_k}(\lambda), f_{\alpha_k}(\lambda))$.

Theorem 12. Suppose α , etc. are proper terms, and $\lambda \in \text{In}_\kappa$.

- a. If $\lambda \notin T_{\alpha_1} \cup T_{\alpha_2}$ then $O(f_{\alpha_1}(\lambda), f_{\alpha_2}(\lambda)) = O(\alpha_1, \alpha_2)$.
- b. If $\lambda \notin T_\alpha$ then $P(f_\alpha(\lambda)) = P(\alpha)$.
- c. If $\lambda \notin T_\alpha$ then $N(f_\alpha(\lambda)) = N(\alpha)$.

Proof. This follows by theorem 10 and lemma 1. □

Theorem 13. Suppose α and β are IV terms, $\alpha = \beta + 1$, $\lambda \in \text{In}_\kappa$, and $\lambda \notin T_\alpha \cup T_\beta$. Then $f_\alpha(\lambda) = f_\beta(\lambda) + 1$.

Proof. In the CNF for α , $\eta_1 = 0$; let α_p be α with the last term σ_1 replaced by τ_1 , or removed entirely if $\sigma_1 = 1$. By theorem 12, $f_\beta(\lambda) = Tf_{\alpha_p}(\lambda)$. It follows that $f_\beta(\lambda) + 1 = Tf_{\alpha_p}(\lambda) + 1 = f_\alpha(\lambda)$. □

6. Set Chains for IV Terms

The reader is assumed to be familiar with [7]. A definition of H_α for an IV term α will be given. A definition of $H_{1\alpha}$ was given in [7] only for a subset of the BV terms. H is a map from subsets of In_κ to subsets of In_κ ; when it is necessary to specify κ , H_α will be denoted $H_{\kappa\alpha}$.

For an IV term α and $X \subseteq \text{In}_\kappa$, say that $\lambda \in H_\alpha(X)$ iff $\lambda \in X$ and $H_\beta(X \cap \lambda)$ is a stationary subset of λ for all IV terms β over λ where $\beta < f_\alpha(\lambda)$.

For $X, Y \subseteq \text{In}_\kappa$ say that $X \subseteq_t Y$ if $\{\lambda \in \text{In}_\kappa : \lambda \in X \text{ and } \lambda \notin Y\}$ is thin; and similarly for \equiv_t .

Theorem 14. If $\alpha' \leq \alpha$ then for any $X \subseteq \text{In}_\kappa$, $H_{\alpha'}(X) \supseteq_t H_\alpha(X)$.

Proof. The proof is by induction on κ . For the basis, κ is the smallest inaccessible cardinal, X is always empty, and the claim is trivial. For arbitrary κ , there is a thin set T such that if $\lambda \in \text{In}_\kappa$ and $\lambda \notin T$ then $f_{\alpha'}(\lambda) \leq f_\alpha(\lambda)$. For such λ , if $H_\beta(X \cap \lambda)$ is stationary for $\beta < f_\alpha(\lambda)$, then by the induction hypothesis $H_\beta(X \cap \lambda)$ is stationary for $\beta < f_{\alpha'}(\lambda)$. □

Lemma 15. Suppose α is an IV term, $X \subseteq \text{In}_\kappa$, $\lambda \in \text{In}_\kappa$, and $\lambda \notin T_\alpha$. Then $H_{\alpha \downarrow \lambda}(X \cap \lambda) = H_\alpha(X) \cap \lambda$.

Proof. Suppose $\mu < \lambda$. Then $\mu \in H_{\alpha \downarrow \lambda}(X \cap \lambda)$ iff $H_\beta(X \cap \mu)$ is stationary for $\beta < f_{\alpha \downarrow \lambda}(\mu)$ iff (by lemma 11) $H_\beta(X \cap \mu)$ is stationary for $\beta < f_\alpha(\mu)$ iff $\mu \in H_\alpha(X)$. □

Theorem 16. *Suppose $\alpha = \beta + 1$, $\lambda \in \text{In}_\kappa$, and $\lambda \notin T_\alpha$. Then $\lambda \in H_\alpha(X)$ iff $\lambda \in H(H_\beta(X))$.*

Proof. $\lambda \in H_\alpha(X)$ iff $\lambda \in X$ and $H_\gamma(X \cap \lambda)$ is stationary for γ with $\gamma < f_\alpha(\lambda)$ iff $\lambda \in H_\beta(X)$ and (A) $H_\gamma(X \cap \lambda)$ is stationary for γ with $\gamma = f_\alpha(\lambda)$. Using lemma 11.b, (A) holds iff $H_{\beta \downarrow \lambda}(X \cap \lambda)$ is stationary, and by lemma 15 this holds iff $H_\beta(X) \cap \lambda$ is stationary. □

Recall from [7] the definition of ρ_R . By theorems 14 and 16, it follows that if $H_\alpha(\text{In}_\kappa)$ is stationary then $\rho_R(H_\alpha(\text{In}_\kappa)) \geq \alpha$. This was proved in [7] only for a subset of the BV terms (with a different definition of H_α ; but see below).

7. A Continuity Property

As noted in [11], explicit ascending sequences can be given for limit ordinals less than Λ_0 , using IV terms. A sequences α_ξ for α may be defined by recursion on α . Case 0 is irrelevant. In case 1, let σ_ξ be as determined by the last node of σ .

Case 2 is divided into subcases as follows. In subcase 2.1, in the CNF for α , $k > 1$. Write $\alpha = \alpha_1 + \alpha_2$. Let α_ξ equal $\alpha_1 + \alpha_{2\xi}$.

In subcase 2.2 $\alpha = \kappa^{+\eta} \cdot \sigma$ where $\sigma \in \text{Lim}$ and $\sigma > 1$. Let α_ξ equal $\kappa^{+\eta} \cdot \sigma_\xi$.

In subcase 2.3 $\alpha = \kappa^{+\eta}$ where $\eta \in \text{Lim}$ and $\eta < \kappa^{+\eta}$. Let α_ξ equal $\kappa^{+\eta_\xi}$.

In subcase 2.4 $\alpha = \kappa^{+\eta}$ where $\eta = \eta^- + 1$. Let α_ξ equal $\kappa^{+\eta^-} \cdot \xi$, where ξ is some particular scheme with ordinal ξ .

To divide case 3 into subcases, write the IV function as $\phi_k(\zeta, \pi_1, \alpha_1, \dots)$. Cases will be denoted XYZ where X is the type of ζ (L for limit, 0, or S for successor), Y is the type of α_1 (L or S), and Z is the type of ξ_1 (L, 0, or S); in some cases Y or Z may be *, denoting any possibility. If α is a successor ordinal write α^- for its predecessor. Sequences α_ξ in each case are as follows.

Case L** : $\alpha_\xi = \phi_k(\zeta_\xi, \pi_1, \alpha_1, \dots)$.

Case 0L* : $\alpha_\xi = \phi_k(0, \pi_1, \alpha_{1\xi}, \dots)$.

Case SL* : $\gamma = \phi_k(\zeta^-, \pi_1, \alpha_1, \dots) + 1$, $\alpha_\xi = \phi_k(\gamma, \pi_1, \alpha_{1\xi}, \dots)$.

Case 0SL : $\alpha_\xi = \phi_k(0, \pi_{1\xi}, \alpha_1^-, \dots)$.

Case SSL : $\gamma = \phi_k(\zeta^-, \pi_1, \alpha_1, \dots) + 1$, $\alpha_\xi = \phi_k(\gamma, \pi_1, \alpha_{1\xi}, \dots)$.

Case 0S0 : $\alpha_0 = \phi_k(0, 0, \alpha_1^-, \dots)$, $\alpha_{n+1} = \phi_k(\alpha_n, 0, \alpha_1^-, \dots)$.

Case SS0 : $\alpha_0 = \phi_k(\zeta^-, 0, \alpha_1, \dots) + 1$, $\alpha_{n+1} = \phi_k(\alpha_n, 0, \alpha_1^-, \dots)$.

Case OSS: $\alpha_0 = \phi_{k+1}(0, \pi_1^-, 1, \pi_1, \alpha_1^-, \dots)$, $\alpha_{n+1} = \phi_{k+1}(0, \pi_1^-, \alpha_n, \pi_1, \alpha_1^-, \dots)$.

Case SSS: $\alpha_0 = \phi_k(\zeta^-, \pi_1, \alpha_1, \dots)$, $\alpha_{n+1} = \phi_{k+1}(0, \pi_1^-, \alpha_n, \pi_1, \alpha_1^-, \dots)$.

Lemma 17. *Suppose α is a term, $\lambda \in In_\kappa$, and $\lambda \notin T_\alpha$. If α is a successor ordinal then $f_\alpha(\lambda)$ is a successor ordinal. If $Cf(\alpha) = \eta$ where $\eta < \kappa$ (whence $\eta < \lambda$) then $Cf(f_\alpha(\lambda)) = \eta$. If $Cf(\alpha) = \kappa$ then $Cf(f_\alpha(\lambda)) = \kappa$. If $Cf(\alpha) = \kappa^+$ then $Cf(f_\alpha(\lambda)) = \lambda^+$.*

Proof. This follows by induction on α , using the ascending sequences given above, and the definition of $f_\alpha(\lambda)$. □

The notation $f \leq_t g$ will be used to denote that $\{\lambda : f(\lambda) > g(\lambda)\}$ is thin; and similarly for $=_t$. It is readily verified that sup and dsup produce least upper bounds in the order \leq_t . See also lemma 1.2 of [2]

Theorem 18. *If $\eta < \kappa$ and α_ξ for $\xi < \eta$ is an ascending sequence with $\alpha = \sup_{\xi < \eta} \alpha_\xi$ then $f_\alpha =_t \sup_{\xi < \eta} f_{\alpha_\xi}$.*

Proof. The proof is by induction on α . By remarks preceding the theorem it suffices to prove the claim for a particular sequence α_ξ . The sequence used will be that given above. Case 0 is irrelevant. Case 1 follows directly. The remaining cases follow inductively, using properties of sup, with subcase 2.4 being irrelevant. □

The theorem can be improved slightly. The set of λ where the equation might not hold is contained in $T_\alpha \cup (\eta + 1) \cup \cup_{\xi < \eta} T_{\alpha_\xi}$.

Theorem 19. *If α_ξ for $\xi < \kappa$ is an ascending sequence with $\alpha = \sup_{\xi < \kappa} \alpha_\xi$ then $f_\alpha =_t dsup_{\xi < \kappa} f_{\alpha_\xi}$.*

Proof. The proof is by induction on α . As in the preceding theorem it suffices to prove the claim for the sequence given above. Case 0 is irrelevant. Case 1 follows directly. The remaining cases follow inductively, using the definition of dsup and properties of sup, with subcase 2.4 and the subcases of case 3 not involving an L being irrelevant. □

Theorem 20. *a. Suppose $\eta < \kappa$, α_ξ is ascending, and $\alpha = \sup_{\xi < \eta} \alpha_\xi$.*

Then $H_\alpha \equiv_t \cap_{\xi < \eta} H_{\alpha_\xi}$.

b. Suppose α_ξ is ascending and $\alpha = \sup_{\xi < \kappa} \alpha_\xi$. Then $H_\alpha \equiv_t$

$Bt_{\xi < \kappa} H_{\alpha_\xi}$.

Proof. For part a, $\lambda \in H_\alpha(X)$ iff $H_\beta(X \cap \lambda)$ is stationary whenever $\beta < f_\alpha(\lambda)$. By theorem 18, except for a thin set of λ , this holds iff, for all $\xi < \eta$, $H_\beta(X \cap \lambda)$ is stationary whenever $\beta < f_{\alpha_\xi}(\lambda)$, iff $\lambda \in H_{\alpha_\xi}(X)$. Part b follows similarly, with “ $\xi < \eta$ ” replaced by “ $\xi < \lambda$ ”. \square

Readily from the definitions, $H_0 = \text{Id}$ (the identity function). By theorem 16, 18, and 19, H_σ as defined here agrees with H^σ as defined in [7]. It may be seen that this continues to hold at κ^+ ; further discussion is omitted here.

8. Enforceability

Let Φ be the sentence $\forall A, C, X, Y, Z$ (“ A is an IV scheme term” \wedge “ C is club” $\wedge X = \text{Inac} \wedge Y = H_A(X) \wedge Z = Y \cap C \Rightarrow Z \neq \emptyset$). Most of the subformulas of the matrix are readily seen to be Δ_0^1 ; discussions will be given for the remaining ones.

There is a Δ_0^1 predicate stating that the class X represents a well-order on a subset of κ . This states that X is a class of ordered pairs, which as a binary relation is transitive and reflexive, and has no descending chains of length ω .

There is a Δ_0^1 predicate stating that the class X represents a scheme for κ . Namely, it represents a pair $\langle \sigma, \phi \rangle$ where σ is represented as above, and ϕ is a function whose domain is the limit points $\alpha < \sigma$, where $\phi(\alpha)$ is a function with domain either an ordinal, or all ordinals, etc.

An IV term may be given as a class coding the sequence of classes $\langle t, \sigma_1, \dots, \sigma_k \rangle$ where t is a hereditarily finite set coding the tree of the term, and $\sigma_1, \dots, \sigma_k$ are (codes for) the schemes at the leaves in order. It follows that the formula “ A is an IV scheme term” is Δ_0^1 .

A function $F : \text{Ord} \mapsto \text{Ord}$ is a class (namely its ordered pairs). Let f_x denote the value of f_σ at stage x of the iteration. The class of triples $\langle x, \lambda, f_x(\lambda) \rangle$ may serve as a witness that $F = f_A$ for a scheme A . The predicate “ W is the witness to $F = f_A$ ” is Δ_0^1 .

The predicate $\alpha = C_{\lambda k}(\eta_1, \sigma_1, \dots, \eta_k, \sigma_k)$, in the values $\lambda, k, \langle \eta_1, \sigma_1, \dots, \eta_k, \sigma_k \rangle$, is first-order definable, indeed no doubt Δ_1^0 . The same is true of the predicate $\alpha = \phi_{\lambda k}(\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k)$, Further discussion is omitted here.

The witness to $F = f_A$ where A is any term is the sequence of classes $\langle F_1, \dots, F_r, W_1, \dots, W_s \rangle$ where r is the number of nodes in the tree of A , s is the number of leaves, for each leaf index i W_i witnesses that $F_{l_i} = f_{\sigma_i}$, and for each interior node index i W_i witnesses that F_{n_i} is obtained by applying the appropriate function per λ to the F_j of the sons. The predicate “ W is the witness to $F = f_A$ ” is Δ_0^1 .

From the foregoing the predicate $F = f_A$ is Σ_1^1 , whence Δ_1^1 .

A term β over $\lambda \in \text{In}_\kappa$ is a set. The predicates “ β is the ordinal of β ” and “ $\gamma = f_\beta(\mu)$ ” are first order definable. The function $\lambda \mapsto S_\lambda$ where $S_\lambda = \{\langle \beta, x, y \rangle : y = H_\beta(x)\}$ may be defined by recursion on λ ; Indeed, $\langle \beta, x, y \rangle \in S_\lambda$ iff $\forall \mu(\mu \in y \text{ iff } \lambda \in x \wedge \forall \beta(\beta \text{ is term over } \mu \wedge \gamma < f_\beta(\mu) \Rightarrow \text{“}H_\gamma(X \cap \mu)\text{ is stationary”}))$.

Now, $Y = H_A(X)$ may be expressed as $\forall \lambda(\lambda \in Y \text{ iff } \lambda \in X \wedge \forall \beta(\beta \text{ is term over } \lambda \wedge \exists \gamma(\langle \lambda, \gamma \rangle \in F \wedge \beta < \gamma \Rightarrow \text{“}H_\beta(X \cap \lambda)\text{ is stationary”}))$. From the foregoing, “ $Y = H_A(X)$ ” is Δ_1^1 .

Say that an inaccessible cardinal κ is Λ_0 -Mahlo iff $\models_{V_\kappa} \Phi$.

Theorem 21. *There is a Π_1^1 sentence Φ such that κ is Λ_0 -Mahlo iff $\models_{V_\kappa} \Phi$.*

Proof. Immediate by remarks above. □

Lemma 22. *Suppose κ is weakly compact. For each scheme σ there is a Π_1^1 formula Φ_σ such that $\models_{V_\kappa} \Phi_\sigma$, and for $\lambda \in \text{In}_\kappa$, if $\models_{V_\lambda} \Phi_\sigma$ then $\sigma \cap V_\lambda$ equals $\sigma \downarrow \lambda$.*

Proof. Φ_σ is constructed by recursion on σ . Various details are omitted, and may be found in [3]. Cases of the recursion are as follows, where Φ has second order parameters α, P . In case 0, Φ is a sentence enforcing that κ is inaccessible. In case 1, Φ_σ is “ σ is a scheme” \wedge “ σ is a successor ordinal” $\wedge \forall \tau(\tau = \sigma^- \Rightarrow \Phi_{\sigma^-}(\tau, P))$. In case 2, Φ_σ is “ σ is a scheme” $\wedge \sigma \in \text{Lim} \wedge \exists \eta(\eta$ is the domain of the ascending sequence of the last node” $\wedge \forall \xi < \eta \forall \tau(\tau = \sigma_\xi \Rightarrow \Phi_{\sigma_\xi}(\tau, P))$). The formulas Φ_{σ_ξ} may be combined into a single formula in a well-known manner. Case 3 is similar. □

Corollary 23. *Suppose κ is weakly compact. For each IV term α there is a Π_1^1 formula Φ_α such that $\models_{V_\kappa} \Phi_\alpha$, and for $\lambda \in \text{In}_\kappa$, if $\models_{V_\lambda} \Phi_\alpha$ then $\alpha \cap V_\lambda$ equals $\alpha \downarrow \lambda$.*

Proof. This follows readily from the lemma and the definition of $\alpha \downarrow \lambda$. □

Theorem 24. *Suppose κ is weakly compact. Then $\models_{V_\kappa} \Phi$.*

Proof. It suffices to show by induction on α that for all α , \models_{V_κ} “ $H_\alpha(\text{Inac})$ is stationary”. Inductively, $\models_{V_\kappa} \forall \beta < \alpha$ “ $H_\beta(\text{Inac})$ is stationary”. By lemma 11 and corollary 23, for a stationary set of λ , $\models_{V_\lambda} \forall \beta < f_\alpha(\lambda)$ “ $H_\beta(\text{Inac})$ is stationary”. \models_{V_κ} “ $H_\alpha(\text{Inac})$ is stationary” follows. □

9. A New Axiom

It seems helpful to define axiom M to be the second order axiom, to be added to NBG, which states that Ord is "...-Mahlo" where "... " is a specification of the length of stationary set chains which have been constructed. As of this paper, axiom M states that "Ord is Λ_0 -Mahlo". The sentence Φ of the previous section gives a more detailed statement.

Only some brief remarks on justifying this axiom will be given here. The reader is assumed to be familiar with section 7 of [6]. Inductively, it may be assumed that $\forall \beta < \alpha$ " $H_\beta(\text{Inac})$ ". Indeed, letting α be a term in some sufficiently large outer universe, successive λ such that $\models_{V_\lambda} \forall \beta < f_\alpha(\lambda)$ " $H_\beta(\text{Inac})$ is stationary" may be collected. The general fact that something that may be repeated may be repeated stationarily often results in a universe in which " $H_\alpha(\text{Inac})$ is stationary" holds.

As usual, more details of this justification would be desirable, but we omit this subject here.

10. Normal Ultrafilters

An application of function chains was given in [5], with further discussion given in [7]. An essential fact is the following.

Theorem 25. *For any measurable cardinal κ , any normal ultrafilter U on κ , and any IV scheme term α over κ , f_α represents α in the ultrapower $Ult_U(V)$.*

Proof. The proof is by induction on α , using Los' theorem. First the claim is proved for schemes. For case 0 the claim is immediate. For case 1, inductively f_τ represents τ , whence $f_{\tau+1}$ represents $\tau+1$. For case 2, inductively f_{σ_ξ} represents σ_ξ , whence $\sup_\xi f_{\sigma_\xi}$ represents $\sup_\xi \sigma_\xi$. Case 3 is similar. For the induction on scheme terms, cases 0 and 1 have already been proved. Cases 2 and 3 follow readily by the definability of the functions C_k and ϕ_k . \square

It is independent whether there are canonical functions of rank κ^+ (see [1]). The foregoing shows that the functions f_α have a weaker property of interest. It would be of interest to obtain a characterization of the property which made no use of ultrafilters.

Suppose κ is a measurable cardinal. Let $o(\kappa)$ denote its Mitchell order. For α an IV scheme term let S_α denote $\{\lambda \in \text{In}_\kappa : o(\lambda) \geq f_\alpha(\lambda)\}$. Recall the definition of $<_R$ from [7].

Theorem 26. *Suppose $\alpha < o(\kappa)$. Then S_α is stationary, and if $\beta < \alpha$ then $S_\beta <_R S_\alpha$.*

Proof. Let U_1 be a normal ultrafilter on κ with $O(U_1) = \alpha$. By lemma 19.34 of [Jech], o represents $o(U_1)$ in $\text{Ult}_{U_1}(V)$. Since by theorem 25 f_α also represents $\alpha = o(U_1)$, $\{\lambda \in \text{In}_\kappa : o(\lambda) = f_\alpha(\lambda)\} \in U_1$. It follows that S_α is stationary.

Suppose $\lambda \in S_{\beta+1}$. Then $o(\lambda) \geq f_{\beta+1}(\lambda)$, so except for a thin set of λ there is a normal ultrafilter U' on λ with $o(U') = f_\beta(\lambda)$. By lemma 11 and theorem 25, except for a thin set of λ , $f_\beta \upharpoonright \lambda$ represents $f_\beta(\lambda)$ in $\text{Ult}_{U'}(V)$, whence it represents $o(U')$. By an argument just given, $\{\mu < \lambda : o(\mu) = f_\alpha(\mu)\} \in U'$, whence $S_\beta \cap \lambda \in U'$. This shows that $S_{\beta+1} \subseteq_t H(S_\beta)$, completing the proof of the theorem. \square

References

- [1] J. Baumgartner, Ineffability properties of cardinals II, In: *Logic, Foundations of Mathematics and Computability Theory* (Ed-s: Butts and Hintikka), Springer, Netherlands (1977), 87-106, **doi:** 10.1007.
- [2] O. Deiser, D. Donder, Canonical functions, non-regular ultrafilters and Ulam's problem on ω_1 , *Journal of Symbolic Logic*, **68** (2003), 713-739, **doi:** 10.2178.
- [3] M. Dowd, Iterating Mahlo's operation, *Int. J. Pure Appl. Math.*, **9**, No. 4 (2003), 469-512.
- [4] M. Dowd, A lower bound on the Mahlo rank of a weakly compact cardinal, *Int. J. Pure Appl. Math.*, **68**, No. 4 (2011), 415-438.
- [5] M. Dowd, Normal ultrafilters and Mahlo rank, *Int. J. Pure Appl. Math.*, **68**, No. 4 (2011), 487-491.
- [6] M. Dowd, Improved results in scheme theory, *Int. J. Pure Appl. Math.*, **76**, No. 2 (2012), 173-190.
- [7] M. Dowd, Scheme terms, *Int. J. Pure Appl. Math.*, **81**, No. 1 (2012), 111-127.
- [8] F. Drake, *Set Theory, An Introduction to Large Cardinals*, North Holland, England (1974).

- [9] K. Schutte, Kennzeichnung von Ordinalzahlen durch rekursiv definierte Funktionen, *Math. Annalen*, **127** (1954), 15-32.
- [10] H. Simmons, Derivatives for ordinal functions and the Schutte brackets, preprint (2004), <http://www.cs.man.ac.uk/~hsimmons/DOCUMENTS/PAPERSandNOTES/Derivatives.dvi>
- [11] O. Veblen, Continuous increasing functions and transfinite ordinals, *Trans. Amer. Math. Soc.*, **9** (1908), 280-292, **doi**: 10.1090.